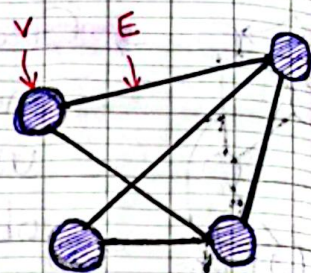


Graph Theory

→ Graphs in Graph Theory

- Collection of nodes and edges
- Denoted by $G = (V, E)$
 - edges (links)
 - vertices (points)
- Size: $n = |V|, m = |E|$



$|V| = 4 \quad |E| = 5$

↳ Vertex / Node

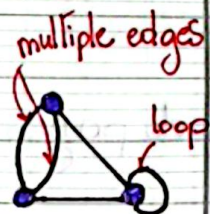
- basic element
- drawn as "node" or "dot"
- denoted by $V(G)$, V , or V_G
- represent object or entity

↳ Edge / Arcs

- Set of 2 elements
- drawn as a line connecting 2 vertices
- called end vertices or endpoints
- denoted by $E(G)$, E or E_G
- represent a relationship

↳ Neighborhood

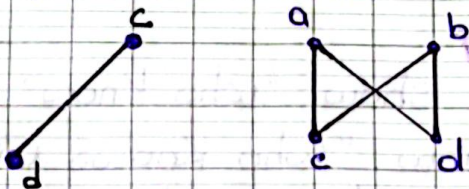
- for any node v the set of nodes it's connected to via an edge is called neighborhood
- denoted by $N(v)$



↳ Simple Graph: no loops, no multiple edges

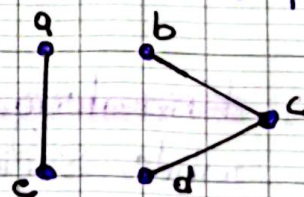
↳ Connected Graphs

There exists at least one path between two vertices



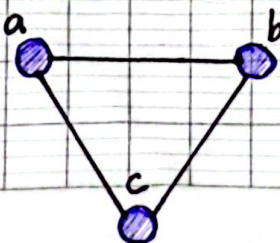
↳ Disconnected

otherwise (separate, unconnect)



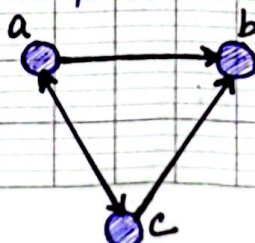
↳ Undirected

2-way, symmetric link



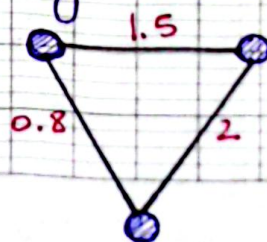
↳ Directed

1-way, asymmetric link, source & target



↳ Weighted

with associated weight (numerical val)



Edge Types

→ Graph Classes

1) Cycle Graph (C_n)

- Single cycle connecting all vertices

$$V = n, E = n$$

2) Complete Graphs

- Every pair of vertices is connected by one edge

$$E = \frac{n(n-1)}{2}$$

3) Bipartite Graphs

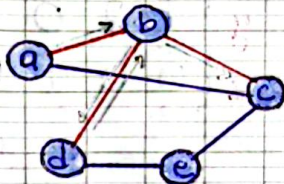
- Vertices divided into two disjoint sets
- No edges within the same set



→ Graph Traversal Concepts

1) Walk

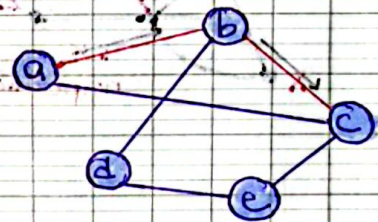
- From v_0 to v_1
- Sequence of vertices and edges
- Vertices and edges can be repeated



ex: $\langle a, b, d, b, c \rangle$

2) Path

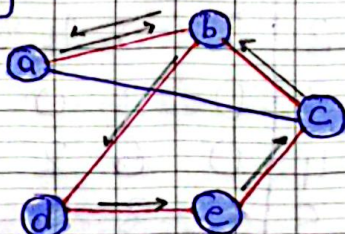
- From v_0 to v_1
- Sequence of vertices
- No vertex is repeated



ex: $\langle a, b, c \rangle$

3) Circuit

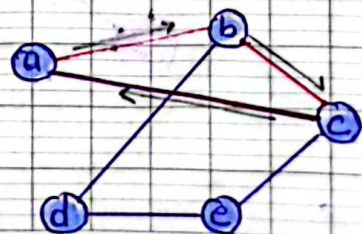
- From v_0 to v_0
- walk that starts and ends at the same vertex
- can repeat vertices and edges



ex: $\langle a, b, d, e, c, b, a \rangle$

4) Cycle

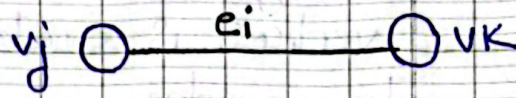
- From v_0 to v_0
- Circuit with no repeated vertices (except start/end)
- Start and end at the same vertex



ex: $\langle a, b, c, a \rangle$

Graph Representation

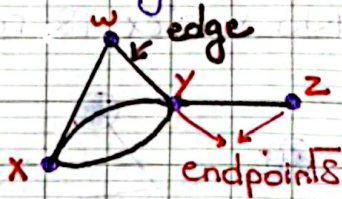
Notations



- vertices (v_j and v_k) said to be **adjacent**
- edge (e_i) said to be **incident** upon v_j
- Degree** of v_k is the nb. of edges incident upon v_k (denoted $d(v_k)$)

3 ways to represent a graph

1) Adjacency matrix $A(G)$

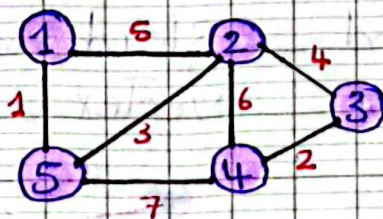


| v_i | w | x | y | z |
|-------|---|---|---|---|
| w | 0 | 1 | 1 | 0 |
| x | 1 | 0 | 2 | 0 |
| y | 1 | 2 | 0 | 1 |
| z | 0 | 0 | 1 | 0 |

$|V| \times |V|$ matrix

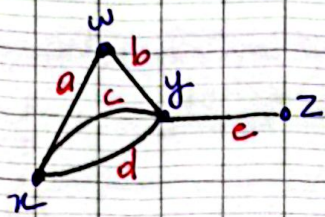
number of edges in G with endpoints $\{v_i, v_j\}$

- Adjacency matrix for a weighted graph
- instead of adding nb. of edges, we add the weight of the edge, if no weight we add 0



| | 1 | 2 | 3 | 4 | 5 |
|---|---|---|---|---|---|
| 1 | 0 | 5 | 0 | 0 | 1 |
| 2 | 5 | 0 | 4 | 6 | 3 |
| 3 | 0 | 4 | 0 | 2 | 0 |
| 4 | 0 | 6 | 2 | 0 | 7 |
| 5 | 1 | 3 | 0 | 7 | 0 |

2) Incidence Matrix



endpoints

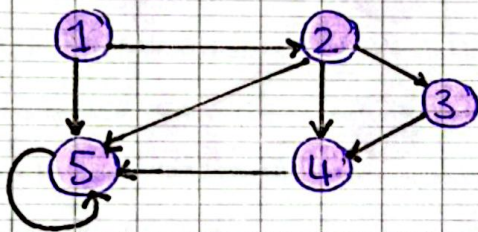
| | a | b | c | d | e ← edges |
|---|---|---|---|---|-----------|
| w | 1 | 1 | 0 | 0 | 0 |
| x | 1 | 0 | 1 | 1 | 0 |
| y | 0 | 1 | 1 | 1 | 1 |
| z | 0 | 0 | 0 | 0 | 1 |

• $|V| \times |E|$ matrix

• 1 if v_i is an endpoint of e_j , 0 otherwise

3) Adjacency List

• Array of $|V|$ elements, one for each vertex in V



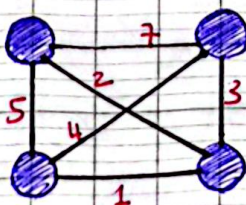
| | | | |
|---|-----|-----|-----|
| 1 | → 2 | → 5 | / |
| 2 | → 5 | → 3 | → 4 |
| 3 | → 4 | / | / |
| 4 | → 5 | / | / |
| 5 | → 5 | / | / |

→ Spanning Tree

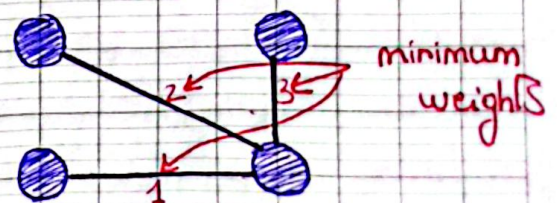
- Subgraph that contains all the vertices and is a tree
- connected, acyclic, has exactly $n-1$ edges
- Graph can have many spanning trees

↳ Minimum Spanning Tree

• Spanning tree with minimum total edge weight



⇒



Complete graph

Minimum Spanning Tree